

Reducing Software Security Risk through an Integrated Approach

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NASA RTOP: Reducing Software Security Risk

NOTE:

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Research Goal

- Reduce security risk to the computing environment by mitigating vulnerabilities in the software development and maintenance life cycles
 - Vulnerability matrix
 - Vulnerabilities exploits and signatures
 - Security Assessment Tools List
 - Property-based testing tool—Tester's Assistant
 - Model-based security specification and verification tool and report



Vulnerability Matrix

- Vulnerability matrix to assist security experts and programmers where best to expend their efforts
 - DOVES database (maintained by UC Davis): http://seclab.cs.ucdavis.edu/projects/
 - Uses the Common Vulnerabilities and Exposures (CVE) Listing (MITRE)
 - http://cve.mitre.org/cve/
 - Contains signatures used to exploit the vulnerability signatures to be used with the Tester's Assistant and the Modeling SPIN Tool



Security Assessment Tools

- Software Security Assessment Instrument
 - > Security assessment tools
 - Description of each tool and its purpose
 - Pros and Cons of each tool
 - Alternate and related tools



Property-Based Testing

- Property-based testing tool Tester's Assistant (Matt Bishop, UC Davis)
 - Perform code slicing on applications for a known set of vulnerabilities
 - ➤ Test for vulnerabilities in code on the system or whenever the computing environment changes
 - ➤ Initially, checks software developed in JAVA
 - The goal is to have the tool check other programming and scripting languages as well (C, C++, Perl, ActiveX, etc.)

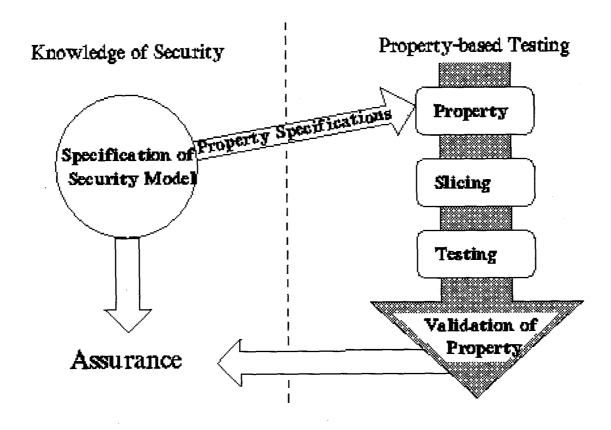


Property-Based Testing (Cont.)

- Compare program actions with specifications
 - Create low-level specifications
 - > Instrument program to check that these hold
 - > Run program under run-time monitor
 - Report violations of specifications



Property-Based Testing (Cont.): How It Works



*Backup Slides provide an example on how this works with the TASPEC



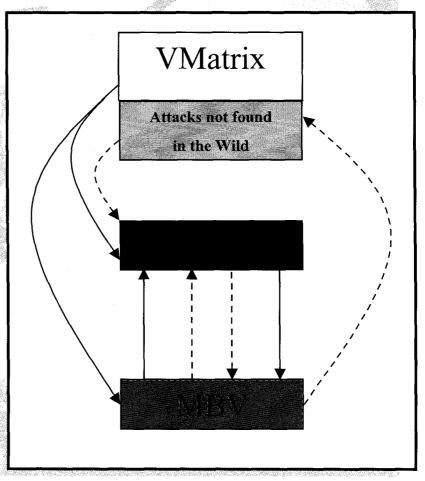
Model-Based Security Specification

- Model-based security specification and verification involves applying formal modeling to the IT security arena
- Verification systems that perform logical verification of temporal properties over models are referred to as model checkers
 - Exhaustive search of a model's corresponding state space
 - Can be used on suitably restricted "partial specifications"



Model Based Verification (MBV) within an Integrated Approach

- Flexible Modeling Framework (FMF)
 - > Compositional Approach
 - > Makes use of SPIN
 - Infers Results from a partial model
- Property Interaction with
 - > Vulnerability (VMatrix)
 - > Property Based Testing (PBT)
- Potentially discovers new vulnerabilities





Real Project Application

- JPL Class A Flight Project
 - ➤ Will test toolset on Flight Mission internet-aware communication software
- IsoWAN & Information Power Grid testbeds
 - Isolated wide-area networks using a modified VPN solution to create a secure, isolated, computing environment
 - Use with high-performance supercomputing collaborative environment



Collaborators

- David Gilliam Principle Investigator

 Network and Computer Security, JPL
- John Powell Research Engineer Quality Assurance, JPL
- John Kelly RTOP Manager
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- Matt Bishop Associate Professor of Computer Science

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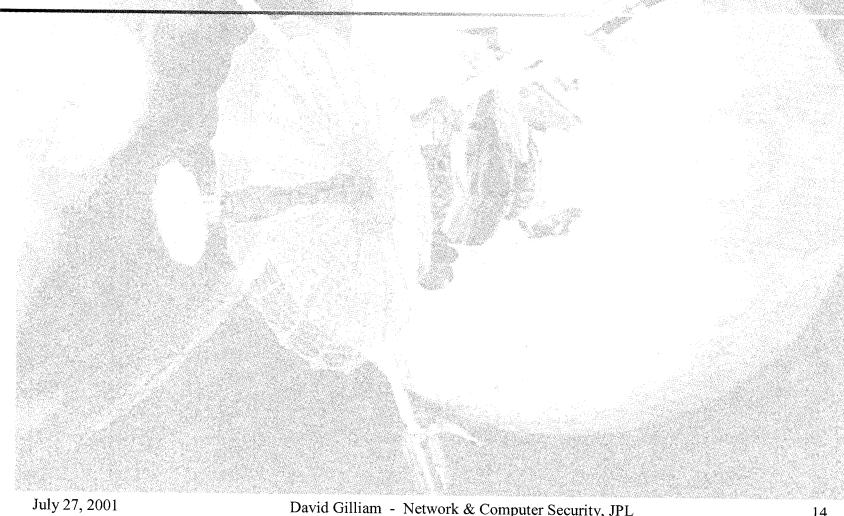
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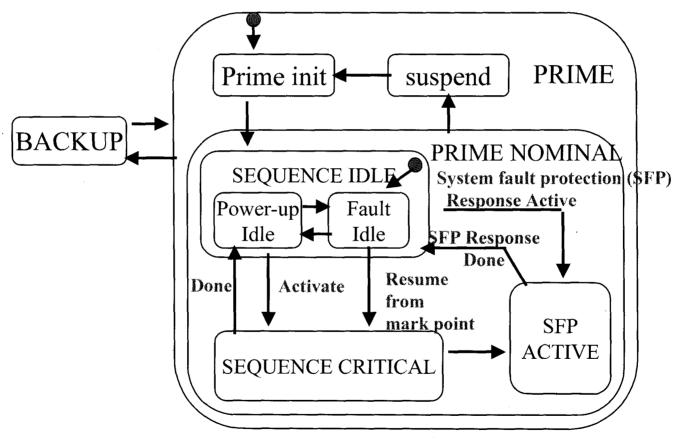
Email: john.d.powell@jpl.nasa.gov



Backup Slides



Model Checking: A Case Study Simplified State Machine for Prime



[&]quot;Validating Requirements for Fault Tolerant Systems Using Model Checking", Schneider, Callahan & Easterbrook, 1998 This Case Study was funded by the NASA Software Program at the NASA IV&V Facility and JPL under a separate task



State Charts

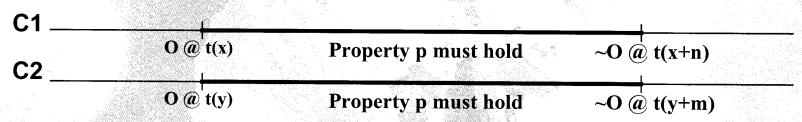
- State Charts are specification notations to define systems
 - Defines the collection of (abstract) variable value pairs at a given point in the system (execution) – referred to as a state
 - Defines the relationships with which the system transitions from one state to the another

The Flexible Modeling Framework (FMF) Approach to MBV

- A Component (c) is some logical unit of process or application behavior
 - A single application often will need to broken into multiple model components
- Combining two components C1 and C2
 - ➢ Model Checking (MC)
 - 1. Non-trivial combination of C1 and C2
 - 2. Searches the Cartesian Product of the sizes of C1 and C2
 - > FMF
 - 1. MC of C1 and C2 individually
 - 2. Combines the State Charts (SC) of C1 and C2
 - 3. Integrates assumptions that follow from 1 above
 - 4. SC traversal or localized MC of appropriate sub-model

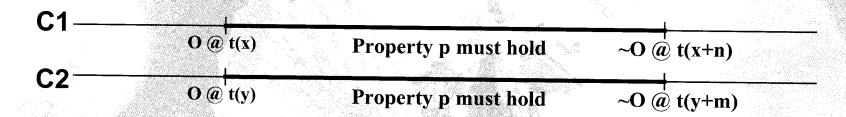


Domain Specifics and FMF



- MC reports p holds for C1 and C2
 - > Assumptions can be made about transitions (T) in C1/C2 SC
 - P holds on T from C1 ^ C2
 - P holds on T from C1 ^ (Unknown in C2)
 - P holds on T from (Unknown in C1) ^ C
- Unify consistent states in the SCs of C1 and C2
 - > Condition: All variables that are known in C1 and C2 agree
- Any path from "O" that does not reach "~O" produces an unknown security result when the combined C1/C2

Combinatorial Network Aware Cases being Addressed



Network Aware (NA) Cases:

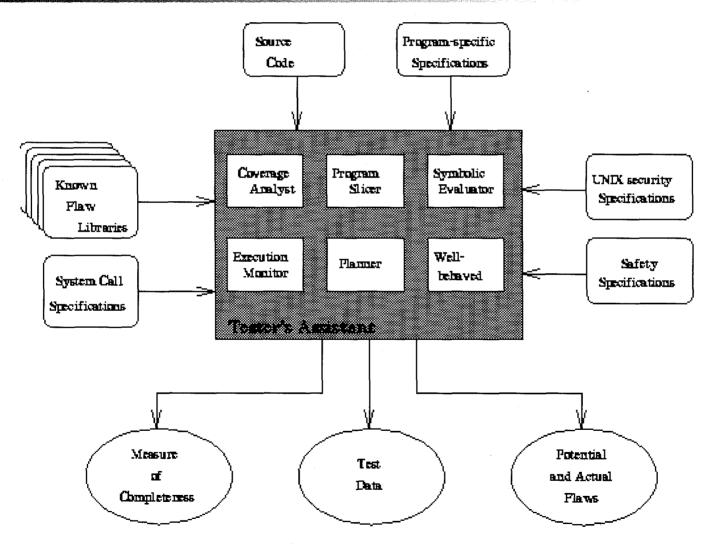
- 1. t(x) = t(y) C1 and C2 are NA simultaneously
- t(x+n) = t(y) C1 ends NA sequence and C2 starts NA sequence simultaneously
- t(x) = t(y+m) C2 ends NA sequence and C1 starts NA sequence simultaneously
- * Sub cases where (n = m) and (n != m) not currently known if this distinction is significant with an abstract model in this domain

Combinatorial Network Aware Cases being Addressed (Cont.)

- The same timing cases seen on the previous slide must be considered in the context of one NA component (C1) and one non-NA component (C2)
 - C1 occurring in a time relation case previously discussed while sharing resources in common may potentially create vulnerabilities.
 - E.g. A NA control application and a printer
 - Non NA components (application pieces) may have been justifiably engineered with little or no consideration of network security issues
 - ➤ A non-NA component may represent a piece of a NA application that does not interact with a network.
 - I.E. t(X+n) < t(y), t(x) > t(y+m)



Property-Based Tester



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David Gilliam - Network & Computer Security, JPL



Property-Based Tester

- TASPEC language definitions
 - > Handle ambiguous specifications and facts
 - > Resetting, non-resetting temporal operators
 - Existential, universal logical operators
- Design Decisions
 - Instrumenter does most work



Tester's Assistant Specifications

 Example: "a user must authenticate himself or herself before acquiring privileges"

```
is password correct? {
    Compare user's password hash to hash stored for that user name If match, set UID to user's uid If no match, set UID to ERROR
}
if privileges granted {
    compare UID to the uid for which privileges are granted if match, all is well if no match, specification violated
}
```



Example C Code

```
if (fgets(stdin, uname, sizeof(uname)-1) == NULL)
    return(FAILED);

typedpwd = getpass("Password: ");
if ((pw = getpwnam(uname)) != NULL){
    hashtp = crypt(pw->pw_passwd, typedpwd);
    if (strcmp(pw->pw_passwd, hashtp) == 0){
        setuid(pw->pw_uid);
        return(SUCCESS);
    }
}
return(FAILED);
```



In TASPEC



Merging

```
if (fgets(stdin, uname, sizeof(uname)-1) == NULL)
    return(FAILED);
                                          user_password(uname, pw->pw_passwd, pw->pw_uid)
typedpwd = getpass("Password: ");
if ((pw = getpwnam (uname)) != NULL){
    hashtp = crypt(pw->pw_passwd, typedpwd);
    if (stremp(pw->pw_passwd, hashtp) == 0){
          setuid(pw->pw uid);
          return(SUCCESS);
                                           user_password(uname, pw->pw_passwd, pw->pw_uid)
return(FAILED);
                                           password_entered(hashtp)
                             user_password(uname, pw->pw_passwd, pw->pw_uid)
                             password_entered(hashtp)
                             equals(pw->pw_passwd, hashtp)
                             authenticated(pw_>pw_uid)
```



Potential Follow-On Work

- Training in use of security assessment tools in the software development and maintenance lifecycle
- Development of re-composable model subcomponents
- Develop capability for easy storage and access of a library of common network security model components and past verification results
- Develop a programmer interface to assist users with generating properties for input into the tools



Potential Follow-On Work (cont.)

- Enhancing and augmenting the toolset
 - > Port the code to run on different operating systems
 - Include additional programming and scripting languages that the Tester's Assistant tool can slice for vulnerabilities
 - Augment the toolset by incorporating or developing additional tools
 - Develop a graphical user interface front-end checklist and decision tree to assist in building the Model to be verified